

# The generatingfunctionology of Clebsch-Gordan coefficients

Robert W. Donley, Jr.  
(CUNY)

June 14, 2019

# Preamble

Wilf, Herbert (1931-2012) , “generatingfunctionology” (book), 1990, 1994, 2006.

This talk: overview of three preprints on arXiv:

- ① (with W. G. Kim), Proceedings of Gestur Olafsson's 65th birthday conference,
- ② Proceedings of CANT 2017/2018,
- ③ arXiv (May 2019), and
- ④ work in progress (magic squares).

Also see pdfs of recent talks at [mathdoctorbob.org](http://mathdoctorbob.org)  
under Group Theory Courses

# Table of Contents:

- ① Introduction
- ② Clebsch-Gordan coefficients for  $SL(2, \mathbb{C})$
- ③ Pascal's Recurrence and Hexagons
- ④ Orthogonality Relations and Frankel's Theorem (1950)
- ⑤ Zeros and 3x3 Magic Squares

## Motivating Question:

Fix  $m, n \geq 0$ .

$V(n)$  be the irr. rep. of  $SL(2, \mathbb{C})$  with highest weight  $n$ .

Then  $\dim_{\mathbb{C}} V(n) = n + 1$ .

### Clebsch-Gordan Decomposition (Gap 2 - Multiplicity One)

$$V(m) \otimes V(n) \cong V(|m - n|) \oplus \cdots \oplus V(m + n - 2) \oplus V(m + n)$$

with no omissions.

Suppose  $m, n$  are even, and non-zero  $\phi_m^0, \phi_n^0$  have weight zero. Then

### Spherical Tensoring (Gap 4)

$$\phi_m^0 \otimes \phi_n^0 \in V(|m - n|) \oplus \cdots \oplus V(m + n - 4) \oplus V(m + n)$$

with no omissions.

(Half Proof: Weyl group element changes parity of  $\phi_m^0$  by  $(-1)^{m/2}$ )

# Clebsch-Gordan Coefficients for $SU(2)$

- ① Classical Invariant Theory:  
A. Clebsch; P. Gordan (Erlangen)
- ② Quantum Mechanics (Weyl, Wigner):  
Coupling of angular momentum and spin  
Squared CGCs as probabilities
- ③ Closed formulas: Wigner, Racah
- ④ Symmetries using Magic Squares: Regge
- ⑤ Vanishing theory: Biedenharn, Louck, Rao, Raynal et al
- ⑥ Simulations in nuclear physics/ chemistry

## Example: $V(2) \otimes V(2) \rightarrow V(4)$

Let  $e, f, h$  be the usual basis for the Lie algebra  $\mathfrak{sl}(2, \mathbb{C})$

Let  $\phi_2$  be a highest weight vector for  $V(2)$

Using the Leibniz rule, a basis for target  $V(4)$  is given by

$$\phi_{2,2} = \phi_2 \otimes \phi_2,$$

$$f^1(\phi_{2,2}) = \phi_2 \otimes f\phi_2 + f\phi_2 \otimes \phi_2,$$

$$f^2(\phi_{2,2}) = \phi_2 \otimes f^2\phi_2 + 2f\phi_2 \otimes f\phi_2 + f^2\phi_2 \otimes \phi_2,$$

$$f^3(\phi_{2,2}) = 3f\phi_2 \otimes f^2\phi_2 + 3f^2\phi_2 \otimes f\phi_2,$$

$$f^4(\phi_{2,2}) = 6f^2\phi_2 \otimes f^2\phi_2.$$

For this talk: Clebsch-Gordan coefficients without normalization for ONB.

That is, we drop unitarity to highlight combinatorial features.

## Example: $V(2) \otimes V(2) \rightarrow V(4)$

- ① array of size  $\dim(first)$  by  $\dim(target) \rightarrow 3 \times 5$ .
- ② columns as coordinate vectors (descending weight),
- ③ Leibniz rule  $\rightarrow$  Pascal's identity (capital-L to right)

$$\begin{bmatrix} 1 & 1 & 1 \\ & 1 & 2 & 3 \\ & & 1 & 3 & 6 \end{bmatrix}$$

Example:  $V(2) \otimes V(2) \rightarrow V(2)$

Choose highest weight vector  $\phi'_{2,2} = \phi_2 \otimes f\phi_2 - f\phi_2 \otimes \phi_2$ .

Using the Leibniz rule, a basis is given by

$$\begin{aligned}\phi'_{2,2} &= \phi_2 \otimes f\phi_2 - f\phi_2 \otimes \phi_2, \\ f^1(\phi'_{2,2}) &= \phi_2 \otimes f^2\phi_2 - f^2\phi_2 \otimes \phi_2, \\ f^2(\phi'_{2,2}) &= f\phi_2 \otimes f^2\phi_2 - f^2\phi_2 \otimes f\phi_2.\end{aligned}$$

## Example: $V(2) \otimes V(2) \rightarrow V(2)$

- ① array of size  $\dim(first)$  by  $\dim(target) \rightarrow 3 \times 3$ .
- ② columns as coordinate vectors (descending weight),
- ③ Leibniz rule  $\rightarrow$  Pascal's identity (capital-L to right)

$$\begin{bmatrix} 1 & 1 \\ -1 & 0 & 1 \\ & -1 & -1 \end{bmatrix}$$

## Cartan Component: $V(m) \otimes V(n) \rightarrow V(m+n)$

Let  $\phi_m$  a highest weight vector for  $V(m)$ .

The same process produces an array of size  $m+1$  by  $m+n+1$ .

First column represents highest weight vector:

$$\phi_{m,n} = \phi_m \otimes \phi_n.$$

Leibniz rule: repeatedly apply  $f$  to highest weight vector.

Clebsch-Gordan coefficients fill a parallelogram section of Pascal's triangle.

Example:  $V(3) \otimes V(4) \rightarrow V(7)$

1	1	1	1	1		
	1	2	3	4	5	
		1	3	6	10	15
			1	4	10	20 35

Decode 7-th column: weight vector of weight  $7 - 2(6) = -5$

- ① need 6 powers of  $f$
- ② row 3 means start with  $f^2$ .

$$f^6(\phi_3 \otimes \phi_4) = 15 f^2\phi_3 \otimes f^4\phi_4 + 20 f^3\phi_3 \otimes f^3\phi_4.$$

## Sub-Cartan Component: $V(m) \otimes V(m) \rightarrow V(2m - 2)$

Let  $\phi_m$  a highest weight vector for  $V(m)$ .

The same process produces an array of size  $m + 1$  by  $2m - 1$ .

First column represents highest weight vector

$$\phi'_{m,m} = \phi_m \otimes f\phi_m - f\phi_m \otimes \phi_m.$$

Leibniz rule: repeatedly apply  $f$  to highest weight vector

Clebsch-Gordan coefficients fill a hexagon of the Catalan “triangle” (after Kirillov-Melnikov).

# Catalan Number Triangle

1	1	1	1	1	1	1	1	1	...
-1	0	1	2	3	4	5	6	7	...
0	-1	-1	0	2	5	9	14	20	...
0	0	-1	-2	-2	0	5	14	28	...
0	0	0	-1	-3	-5	-5	0	14	...
0	0	0	0	-1	-4	-9	-14	-14	...
...	...	...	...	...	...	...	...	...	...

Catalan Numbers (to the right of the diagonal zeros)

1, 1, 2, 5, 14, 42, 132, 429, ...

## Catalan Number Triangle

1	1	1	1	1	1	1	1	1	...
-1	0	1	2	3	4	5	6	7	...
0	-1	-1	0	2	5	9	14	20	...
0	0	-1	-2	-2	0	5	14	28	...
0	0	0	-1	-3	-5	-5	0	14	...
0	0	0	0	-1	-4	-9	-14	-14	...
...	...	...	...	...	...	...	...	...	...

*First column  $\rightarrow p(x) = 1 - x$*

$$(1 - x)(1 + x) = 1 + 0x - x^2$$

$$(1 - x)(1 + x)^2 = 1 + x - x^2 - x^3$$

$$(1 - x)(1 + x)^3 = 1 + 2x + 0x^2 - 2x^3 - x^4$$

$$(1 - x)(1 + x)^4 = 1 + 3x + 2x^2 - 2x^3 - 3x^4 - x^5$$

## Example:

$$V(3) \otimes V(3) \rightarrow V(4) : \quad 4 \times 5$$

$$3 \begin{bmatrix} 1 & 1 & 1 & \\ -1 & 0 & 1 & 2 \\ & -1 & -1 & 0 & 2 \\ & & -1 & -2 & -2 \end{bmatrix}$$

$$V(4) \otimes V(4) \rightarrow V(6) : \quad 5 \times 7$$

$$4 \begin{bmatrix} 1 & 1 & 1 & 1 & \\ -1 & 0 & 1 & 2 & 3 \\ & -1 & -1 & 0 & 2 & 5 \\ & & -1 & -2 & -2 & 0 & 5 \\ & & & -1 & -3 & -5 & -5 \end{bmatrix}$$

$$\text{Hexagon: } V(m) \otimes V(n) \rightarrow V(m+n-2k)$$

Given non-negative  $m, n, k$  with  $0 \leq k \leq \min(m, n)$

- ① Initialize matrix of size  $\dim(\text{first})$  by  $\dim(\text{target})$ ,
- ② Initialize left-most column with highest weight condition,
- ③ Pascal's Recurrence (Capital-L),
- ④ Remove upper-right triangle for symmetry.

Columns are coordinate vectors for weight vectors of target space  
(descending weights)

## Generating Function for Highest Weight Coefficients

The highest weight vector  $\phi_{m,n}^k$  for  $V(m+n-2k)$  is a linear combination of terms

$$\phi_{m,n}^k = \sum_{i+j=k} c(i,j) f^i \phi_m \otimes f^j \phi_n.$$

The coefficient of

$$f^i \phi_m \otimes f^j \phi_n \quad \text{in} \quad \phi_{m,n}^k$$

is the coefficient of  $x^j y^i$  in the expansion of

$$\frac{1}{(1-x)^{m-k+1}} \cdot \frac{1}{(1+y)^{n-k+1}}.$$

Specifically, in row  $i+1$

$$c_{m,n,k}(i, k-i) = (-1)^i \binom{m-i}{k-i} \binom{n-k+i}{i}$$

## Generating Function for CG Coefficients

The vector  $f^t \phi_{m,n}^k$  for  $V(m+n-2k)$  is a linear combination of terms

$$f^t \phi_{m,n}^k = \sum_{i+j=t+k} c(i,j) f^i \phi_m \otimes f^j \phi_n.$$

The coefficient of

$$f^i \phi_m \otimes f^j \phi_n \quad \text{in} \quad f^t \phi_{m,n}^k$$

is the coefficient of  $x^j y^i$  in the expansion of

$$(x+y)^t \cdot \frac{1}{(1-x)^{m-k+1}} \cdot \frac{1}{(1+y)^{n-k+1}}.$$

Specifically, in row  $i+1$  and column  $i+j-k+1$ ,

$$c_{m,n,k}(i,j) = \sum_{s=0}^k (-1)^s \binom{i+j-k}{i-s} \binom{m-s}{k-s} \binom{n-k+s}{s}$$

# Wigner's Formula (1931)

For instance, Vilenkin's book, "Special Functions..." (1965).

$$C_{m,n,k}(i,j) = \sqrt{\frac{(m+n-2k+1)(m+n-i-j-k)!(i+j-k)!i!j!k!}{(m-i)!(n-j)!(m-k)!(n-k)!(m+n-k+1)!}} \cdot S$$

where

$$S = \sum_{s=0}^k (-1)^{i-s} \frac{(m-i+s)!(n-k+i-s)!}{s!(i+j-k-s)!(i-s)!(k-i+s)!}$$

Traditionally studied through representation theory, Jacobi polynomials and hypergeometric series of type  ${}_3F_2$  (terminating)

# Wigner's Formula (1931)

With binomial coefficients and shift  $s \rightarrow i - s$ ,

$$C_{m,n,k}(i,j) = \sqrt{\frac{(m+n-2k+1) \binom{m+n-k}{m-i, n-j, *} }{(m+n-k+1) \binom{m+n-k}{m-k, n-k, k} \binom{m+n-k}{i, j, *} }} \cdot S$$

where

$$S = \sum_{s=0}^k (-1)^s \binom{i+j-k}{i-s} \binom{m-s}{k-s} \binom{n-k+s}{s}$$

# Techniques:

Traditional:

- ① ONB and normalizations
- ② Special functions and hypergeometric series
- ③ Weighted recurrences

Characteristic zero, assuming highest weight rep's:

- ① No normalizations → Integer coefficients
- ② Uniform generating function
- ③ Pascal's recurrence (Calculus of finite differences)

See quantum CGC theory (Kirillov-Reshetikhin (1988) *et al*)

# Combinatorial Toolbox:

Uniform generating function →

- ① Recurrence relations (three and four term)
- ② Explicit calculations of summations using Dixon's formula
- ③ Elementary algorithm for programming (MAPLE and EXCEL)
- ④ **Orthogonality relations**
- ⑤ **Regge symmetries (magic squares)**

# Topic: Orthogonality Relations

Change between orthonormal bases  $\rightarrow$  unitary matrix

$$\sum_{j=1}^n c_{i,j} \overline{c_{i',j}} = \delta_{i,i'}$$

Non-degenerate invariant bilinear forms:

if  $X$  is in  $\mathfrak{sl}(2, \mathbb{C})$ , then

$$\langle Xu, v \rangle = -\langle u, Xv \rangle.$$

For example,

$$\langle f^i \phi_m, f^{m-i} \phi_m \rangle = (-1)^i \langle \phi_m, f^m \phi_m \rangle \neq 0$$

and

$$f^i \phi_m \otimes f^j \phi_n \quad \text{pairs with} \quad f^{m-i} \phi_m \otimes f^{n-j} \phi_n.$$

## Orthogonality Relations (Non-vanishing Part)

The sum

$$\sum_{i+j=s} c_{m,n,k}(i,j) c_{m,n,k}(m-i, n-j)$$

is independent of  $s$ .

## Example:

$V(3) \otimes V(3) \rightarrow V(4)$  :

$$\begin{bmatrix} 1 \\ -1 \\ 0 \\ 0 \end{bmatrix} \begin{bmatrix} 1 & 1 & * \\ 0 & 1 & 2 \\ -1 & -1 & 0 \\ 0 & -1 & -2 \end{bmatrix} \begin{bmatrix} * \\ * \\ 2 \\ -2 \end{bmatrix},$$

$$\begin{bmatrix} 1 & 1 & 1 & * \\ -1 & 0 & 1 & * \\ 0 & -1 & -1 & 2 \\ 0 & 0 & -1 & 0 \end{bmatrix} \begin{bmatrix} * \\ * \\ 2 \\ -2 \end{bmatrix}$$

$$(1, -1, 0, 0) \cdot (-2, 2, *, *) = -4$$

$$(1, 0, -1, 0) \cdot (-2, 0, 2, *) = -4$$

$$(1, 1, -1, -1) \cdot (-1, -1, 1, 1) = -4$$

# Generalized Binomial Transform

$$\{a_i\}_{i=0}^{\infty} \rightarrow p(x) = \sum_{i=0}^{\infty} a_i x^i$$

## Generalized Binomial Transform

$$B^n a_k = a_{k,n} = \sum_{i=0}^n \binom{n}{i} a_{k-i}$$

$$\sum_{k=0}^{\infty} B^n a_k x^k = (1+x)^n \sum_{i=0}^{\infty} a_i x^i$$

Extend to all  $n$  using

$$(1+x)^{-n} = \sum_{i=0}^{\infty} (-1)^i \binom{n+i-1}{i} x^i$$

# Pascal's Recurrence

$$n = 1 : B^1 a_k = a_{k-1} + a_k$$

$$n = 2 : B^2 a_k = a_{k-2} + 2a_{k-1} + a_k$$

$$n = 3 : B^3 a_k = a_{k-3} + 3a_{k-2} + 3a_{k-1} + a_k$$

## Pascal's Recurrence (Capital L)

$$B^{n+1} a_k = B^n a_k + B^n a_{k-1}$$

## Pascal's Identity

$$\binom{n+1}{k} = \binom{n}{k} + \binom{n}{k-1}$$

# Matrix Implementation (Binomial Array)

- ① Fourth quadrant matrix
- ②  $\{a_i\}$  down first column,  $a_0$  along first row
- ③ Pascal's recurrence: Capital-L summation

$$\begin{bmatrix} a_0 & a_0 & a_0 & a_0 & a_0 & \cdots \\ a_1 & a_0 + a_1 & 2a_0 + a_1 & 3a_0 + a_1 & 4a_0 + a_1 & \cdots \\ a_2 & a_1 + a_2 & a_0 + 2a_1 + a_2 & 3a_0 + 3a_1 + a_2 & 6a_0 + 4a_1 + a_2 & \cdots \\ a_3 & a_2 + a_3 & a_1 + 2a_2 + a_3 & \cdots & \cdots & \cdots \\ \cdots & \cdots & \cdots & \cdots & \cdots & \cdots \end{bmatrix}$$

$B^n a_k$ : row  $k + 1$ , column  $n + 1$  (first row and column indexed to 0)

# Catalan Number Triangle

$$a_i = (1, -1, 0, 0, \dots), \quad B^n a_i = \binom{n}{i} - \binom{n}{i-1}$$

1	1	1	1	1	1	1	1	1	...
-1	0	1	2	3	4	5	6	7	...
0	-1	-1	0	2	5	9	14	20	...
0	0	-1	-2	-2	0	5	14	28	...
0	0	0	-1	-3	-5	-5	0	14	...
0	0	0	0	-1	-4	-9	-14	-14	...
...	...	...	...	...	...	...	...	...	...

Right of zeros:

$$B^{2n} a_n = \frac{1}{n+1} \binom{2n}{n} = C_n$$

# Discrete Convolution

## Discrete convolution

If  $a_i$  and  $b_j$  are sequences, we define a new sequence, the **discrete convolution** (or Cauchy product) by

$$(a * b)_n = \sum_{i+j=n} a_i b_j = \sum_{i=0}^n a_i b_{n-i}.$$

Alternatively,  $(a * b)_n$  is the coefficient  $c_n$  in the power series product

$$\sum_{i=0}^{\infty} c_i x^i = \left( \sum_{j=0}^{\infty} a_j x^j \right) \left( \sum_{k=0}^{\infty} b_k x^k \right).$$

## Example: Catalan Numbers

$$C_0 = 1, \quad C_{n+1} = \sum_{i=0}^n C_i C_{n-i}$$

Implement as iterated discrete convolution:

$$C_1 = 1 \cdot 1 = 1,$$

$$C_2 = (1, 1) \cdot (1, 1) = 2,$$

$$C_3 = (1, 1, 2) \cdot (2, 1, 1) = 5,$$

$$C_4 = (1, 1, 2, 5) \cdot (5, 2, 1, 1) = 14,$$

$$C_5 = (1, 1, 2, 5, 14) \cdot (14, 5, 2, 1, 1) = 42, \dots$$

Closed formula:

$$C_n = \frac{1}{n+1} \binom{2n}{n}$$

# Frankel's Theorem (1950)

## Theorem (Dwyer, Frankel)

If  $a_i$  and  $b_j$  are sequences, then, for all  $n$  in  $\mathbb{Z}$ ,

$$(a * b)_k = (B^n a * B^{-n} b)_k.$$

## Interpretation if $a_i = b_i$

- ① construct array for  $B^n a_k$ ,
- ② section off any rectangle from the top row,
- ③ the convolution of the left and right hand columns is unchanged under telescoping.

## Application: See Shapiro (1976)

### Theorem (D.)

The length-squared of the  $n$ -th column of the Catalan triangle is  $2C_n$ .

Proof:

$$\left[ \begin{array}{ccccccccc} 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & \dots \\ -1 & 0 & 1 & 2 & 3 & 4 & 5 & 6 & \dots \\ 0 & -1 & -1 & 0 & 2 & 5 & 9 & 14 & \dots \\ 0 & 0 & -1 & -2 & -2 & 0 & 5 & 14 & \dots \\ 0 & 0 & 0 & -1 & -3 & -5 & -5 & 0 & \dots \\ 0 & 0 & 0 & 0 & -1 & -4 & -9 & -14 & \dots \\ \dots & \dots \end{array} \right]$$



## Proof of Frankel's Theorem:

Theorem (Dwyer, Frankel)

If  $a_i$  and  $b_j$  are sequences, then, for all  $n$  in  $\mathbb{Z}$ ,

$$(a * b)_k = (B^n a * B^{-n} b)_k.$$

Suppose  $p(x)$  and  $q(x)$  correspond to  $a_i$  and  $b_j$ , respectively.

Define  $[x^k]p(x) = a_k$ . Then  $(a * b)_k = [x^k](p(x)q(x))$ .

$$\begin{aligned}(B^n a * B^{-n} b)_k &= [x^k]((1+x)^n p(x)(1+x)^{-n} q(x)) \\&= [x^k](p(x)q(x)) \\&= (a * b)_k. \quad \square\end{aligned}$$

# Topic: Magic Squares and Zeros of CGCs

Regge (1958):

domain space for CGCs corresponds precisely to 3x3 magic squares

$$V(m) \otimes V(n) \rightarrow V(m+n-2k)$$

$$M = \begin{bmatrix} n-k & m-k & k \\ i & j & * \\ m-i & n-j & i+j-k \end{bmatrix} \mapsto c_{m,n,k}(i,j)$$

Magic number:  $J = m + n - k$

Recall:  $f^i \phi_m \otimes f^j \phi_n$  pairs with  $f^{m-i} \phi_m \otimes f^{n-j} \phi_n.$

# $\mathbb{M}_3 = 3 \times 3$ weakly semi-magic squares

$$V(b+k) \otimes V(a+k) \rightarrow V(a+b)$$

$$M = \begin{bmatrix} a & b & k \\ r & * & * \\ * & * & c \end{bmatrix}$$

- ① all entries are nonnegative integers,
- ② all line sums along rows and columns are equal, and
- ③ this sum, the **magic number**, equals  $J = a + b + k$ .

Also called **integer “doubly-stochastic” matrices**.

## Clebsch-Gordan coefficients/ function

$$C : \mathbb{M}_3 \rightarrow \mathbb{Z}$$

$$M = \begin{bmatrix} a & b & k \\ r & m & * \\ * & * & c \end{bmatrix} \mapsto$$

$$C(M) = \sum_{l=0}^k (-1)^l \binom{c}{r-l} \binom{b+k-l}{b} \binom{a+l}{a}$$

$C(M)$  is the coefficient of

$$x^m y^r$$

in the power series expansion of

$$\frac{(x+y)^c}{(1-x)^{b+1}(1+y)^{a+1}}$$

# Determinantal Symmetries

For  $3 \times 3$  matrices, let  $G$  be the group of determinantal symmetries;  
that is,

- ①  $G$  is generated by row switches, column switches, and transpose,
- ② every element  $g$  of  $G$  may be expressed uniquely as

$$g = R(\sigma)C(\tau)T^\epsilon \quad \text{with} \quad \sigma, \tau \in S_3,$$

and

- ③  $|G| = 72$ .

These symmetries preserve

- ① the semi-magic square property,
- ② the magic number  $J$ , and
- ③ the zero locus for CGCs. (Regge, 1958)

## Classification of Zeros:

Open problem: Classify the zeros of  $C(M)$ .

(Biedenharn, Brudno, Louck, K. S. Rao, etc.)

How to organize  $\begin{bmatrix} a & b & k \\ r & * & * \\ * & * & c \end{bmatrix}$  as a subset of  $\mathbb{N}^5$ ?

Three steps:

- ① Fix a magic number  $J \geq 0$ ,
- ② partition by top-lines  $\rightarrow$  equilateral triangle of size  $J + 1$ , and
- ③ each top-line corresponds to a hexagon (as seen before).

That is, the data attached to

$$V(b+k) \otimes V(a+k) \rightarrow V(a+b)$$

is indexed by all magic squares with top line

$$(a, b, k).$$

## Example: Hexagon without CGCs

$J = 4$ , top line:  $(1, 2, 1)$   $\rightarrow$  10 squares

$$V(3) \otimes V(2) \rightarrow V(3)$$

$$\begin{bmatrix} 1 & 2 & 1 \\ r & * & * \\ * & * & c \end{bmatrix} \mapsto \begin{bmatrix} * & * \\ * & * & * \\ * & \boxed{*} & * \\ * & * & * \end{bmatrix}$$

Box corresponds to  $(r, c) = (2, 2)$ :

$$\begin{bmatrix} 1 & 2 & 1 \\ 2 & 1 & 1 \\ 1 & 1 & 2 \end{bmatrix}$$

Note: No magic square for  $(r, c) = (2, 0)$  if  $J = 4$ .

## Example: Triangle of Top Lines

$$\mathbf{J} = 4$$

$(0, 0, 4), (0, 1, 3), (0, 2, 2), (0, 3, 1), (0, 4, 0)$

$(1, 0, 3), (1, 1, 2), \boxed{(1, 2, 1)}, (1, 3, 0)$

$(2, 0, 2), (2, 1, 1), (2, 2, 0)$

$(3, 0, 1), (3, 1, 0)$

$(4, 0, 0)$

Top line entries: (rows above, spaces to left, spaces to right)

## Example: Triangle with Magic Square Counts

$$\mathbf{J} = 4$$

5, 8, 9, 8, 5

8, 10, 10, 8

9, 10, 9

8, 8

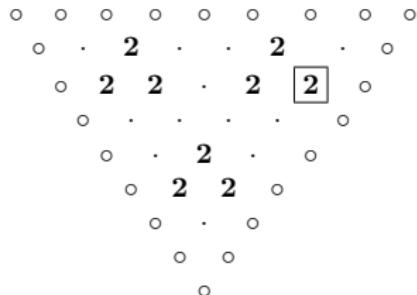
5

Total squares = 120

## Example: $J = 8$

#	Magic Squares
9	16 21 24 25 24 21 16 9
16	22 26 28 28 26 22 16
21	26 29 30 29 [26] 21
24	28 30 30 28 24
25	28 29 28 25
24	26 26 24
21	22 21
16	16
9	

## # CGC Zeros



Top line for box:     2 above, 5 to the left, 1 to the right  $\mapsto (2, 5, 1)$

Example:  $J = 8 : V(6) \otimes V(3) \rightarrow V(7)$

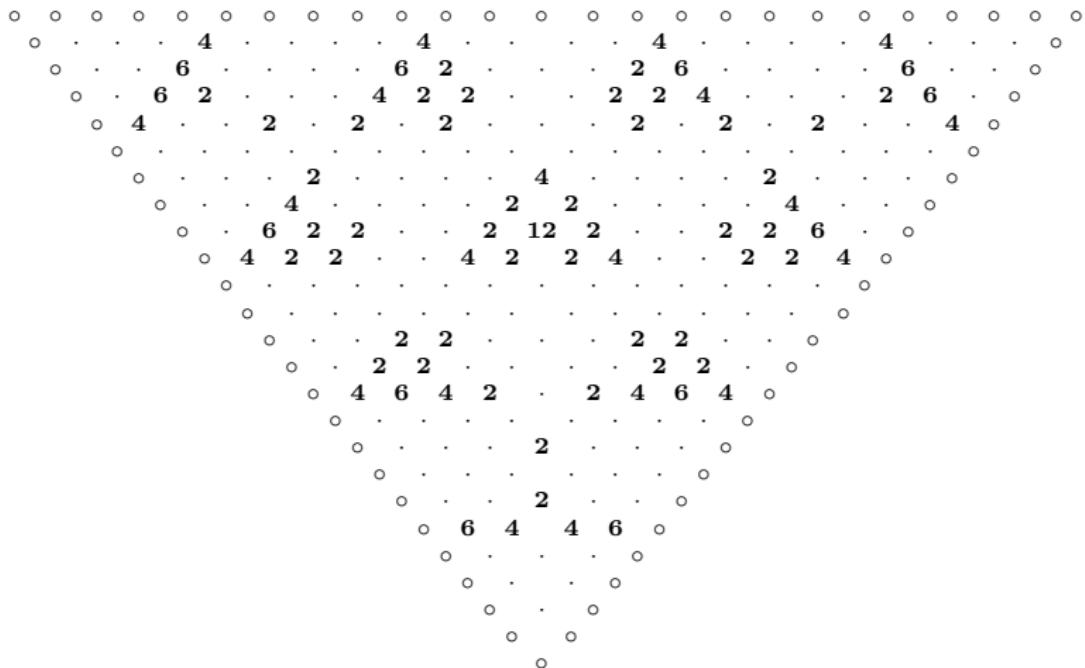
Top line (2, 5, 1) : 26 squares, 2 zeros

$$\begin{bmatrix} 6 & 6 & 6 \\ -3 & 3 & 9 & 15 \\ -3 & \boxed{0} & 9 & 24 \\ -3 & -3 & 6 & 30 \\ -3 & -6 & \boxed{0} & 30 \\ -3 & -9 & -9 & 21 \\ -3 & -12 & -21 \end{bmatrix}$$

Zeros:  $(r, c) = (2, 2), (4, 5)$

$$\begin{bmatrix} 2 & 5 & 1 \\ 2 & 1 & 5 \\ 4 & 2 & 2 \end{bmatrix}, \quad \begin{bmatrix} 2 & 5 & 1 \\ 4 & 2 & 2 \\ 2 & 1 & 5 \end{bmatrix}$$

Example:  $J = 24$  : 252 Zeros, 6 Orbits



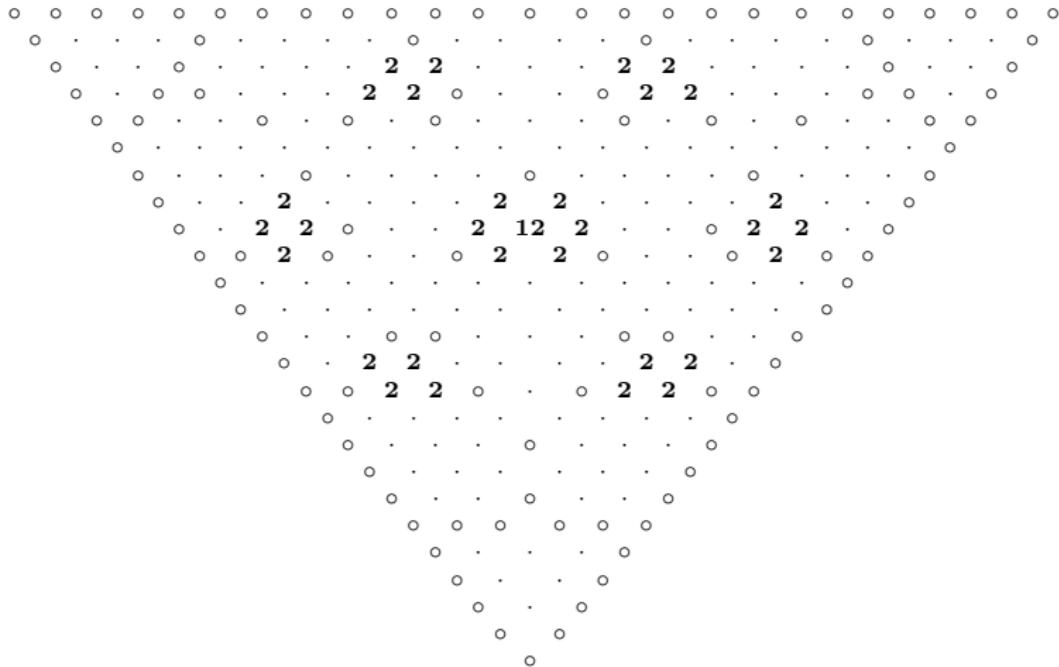
Example:  $J = 24 : V(16) \otimes V(16) \rightarrow V(16)$

Top Line (8,8,8), 12 zeros

Zeros are orbits under column switches, R23 switch of

$$M = \begin{bmatrix} 8 & 8 & 8 \\ 2 & 9 & 13 \\ 14 & 7 & 3 \end{bmatrix}$$

# Example: $J = 24$ : Full orbit of 72 Zeros



*Thank you!*