

# Clebsch-Gordan Coefficients in Characteristic Zero

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# Acknowledgments

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**m=3, n=4, k=0**

$$(m + 1) \times (m + n - 2k + 1) = 4 \times 8$$

$k = 0$ : Pascal's Triangle (truncate at right corner for symmetry)

$$\begin{bmatrix} 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 \\ 0 & 1 & 2 & 3 & 4 & 5 & 0 & 0 \\ 0 & 0 & 1 & 3 & 6 & 10 & 15 & 0 \\ 0 & 0 & 0 & 1 & 4 & 10 & 20 & 35 \end{bmatrix}$$

**m=3, n=4, k=1:**

$$(m + 1) \times (m + n - 2k + 1) = 4 \times 6$$

$k = 1$  : Set left-most column -  $k + 1 = 2$  nonzero entries,  
Pascal's Recurrence to Right

$$\begin{bmatrix} 3 & 3 & 3 & 3 & 0 & 0 \\ -4 & -1 & 2 & 5 & 8 & 0 \\ 0 & -4 & -5 & -3 & 2 & 10 \\ 0 & 0 & -4 & -9 & -12 & -10 \end{bmatrix}$$

$m=3, n=4, k=2$

$$(m + 1) \times (m + n - 2k + 1) = 4 \times 4$$

$k = 2$  : Set left-most column -  $k + 1 = 3$  nonzero entries,  
Pascal's Recurrence to Right

$$\begin{bmatrix} 3 & 3 & 3 & 0 \\ -6 & -3 & 0 & 3 \\ 6 & 0 & -3 & -3 \\ 0 & 6 & 6 & 3 \end{bmatrix}$$

**m=3, n=4, k=3**

$$(m + 1) \times (m + n - 2k + 1) = 4 \times 2$$

$k = 3$  : Set left-most column -  $k + 1 = 4$  nonzero entries,  
Pascal's Recurrence to Right

$$\begin{bmatrix} 1 & 1 \\ -2 & -1 \\ 3 & 1 \\ -4 & -1 \end{bmatrix}$$

## First $k + 1$ Entries Of The Left-Most Column

$$(-1)^s \binom{m-s}{k-s} \binom{n-k+s}{s}, \quad 0 \leq s \leq k$$

$$k = 1: \quad \left( \binom{3}{1} \binom{3}{0}, -\binom{2}{0} \binom{4}{1} \right) \rightarrow (3, -4, 0, 0)$$

$$k = 2: \quad \left( \binom{3}{2} \binom{2}{0}, -\binom{2}{1} \binom{3}{1}, \binom{1}{0} \binom{4}{2} \right) \rightarrow (3, -6, 6, 0)$$

$$k = 3: \quad \left( \binom{3}{3} \binom{1}{0}, -\binom{2}{2} \binom{2}{1}, \binom{1}{1} \binom{3}{2}, -\binom{0}{0} \binom{4}{3} \right) \\ \rightarrow (1, -2, 3, -4)$$

# The Algorithm

Given non-negative  $m$ ,  $n$ ,  $k$  with  $0 \leq k \leq \min(m, n)$

1. Initialize matrix of size  $(m + 1)$  by  $(m + n - 2k + 1)$ ,
2. Initialize left-most column with binomial coefficient pattern,
3. Pascal's Recurrence to the Right,
4. Zero out upper-right column for symmetry.

Columns are coordinate vectors associated to a tensor product of vector spaces

# Main Result

Fix  $m, n, k$  as before. The  $(i + 1, i + j - k + 1)$  entry of the corresponding matrix in the algorithm equals

$$\begin{aligned}c_{m,n,k}(i,j) &= \sum_{s=0}^k (-1)^s \binom{i+j-k}{i-s} \binom{m-s}{k-s} \binom{n-k+s}{s} \\ &= \sum_{s=0}^k (-1)^s \binom{i+j-k}{i-s} \binom{m-i}{k-s} \binom{n-j}{s}\end{aligned}$$

Pascal's Recurrence:

$$c_{m,n,k}(i,j) = c_{m,n,k}(i-1,j) + c_{m,n,k}(i,j-1)$$

# Generating Function

$$c_{m,n,k}(i,j) = \sum_{s=0}^k (-1)^s \binom{i+j-k}{i-s} \binom{m-s}{k-s} \binom{n-k+s}{s}$$

is the coefficient of  $x^j y^i$  in the expansion of

$$(x+y)^{i+j-k} \cdot \frac{1}{(1-x)^{m-k+1}} \cdot \frac{1}{(1+y)^{n-k+1}}$$

This yields a second Pascal-type recurrence

$$c_{m,n,k}(i,j) = c_{m-1,n,k}(i-1,j) + c_{m,n-1,k}(i,j-1)$$

# Clebsch-Gordan Coefficients

$$C_{m,n,k}(i,j) = \frac{\binom{m+n-k}{m-i, n-j, *}}{\binom{m+n-k}{i, j, *}} C_{m,n,k}(i,j) \in \mathbb{Q}$$

where

$$D(m, n, k) = \binom{m+n-k+1}{k} \binom{m+n-2k}{m-k}$$

# Clebsch-Gordan Coefficients

1. Classical Invariant Theory
2. Tensor products of finite-dimensional representations of  $SL(2, R)$ ,  $SU(2)$ ,  $SL(2, C)$
3. Quantum Theory of Angular Momentum
4. Multiplication Formula for Jacobi Polynomials and other Special Functions
5. Modular forms: Rankin-Cohen brackets

# Clebsch-Gordan Decomposition

$V(m)$  = irreducible finite-dimensional representation for  $SU(2)$   
Highest weight =  $m$ , dimension =  $m + 1$

$0 \leq m \leq n$ :

$$\begin{aligned} V(m) \otimes V(n) &= \bigoplus_{k=0}^m V(m+n-2k) \\ &= V(n-m) \oplus V(n-m+2) \oplus \cdots \oplus V(n+m) \end{aligned}$$

Example:  $k = 3, 2, 1, 0$

$$V(3) \otimes V(4) = V(1) \oplus V(3) \oplus V(5) \oplus V(7)$$

$$4 \cdot 5 = 2 + 4 + 6 + 8$$

# Clebsch-Gordan Decomposition - Vector Level

$V(m)$ : basis  $\{u_i\}$

$V(n)$ : basis  $\{v_j\}$

Natural basis for  $V(m) \otimes V(n)$ :  $\{u_i \otimes v_j\}$

Bases for each  $V(m + n - 2k)$ :  $\{w_p^k\}$

$$w_p^k = \sum_{i+j} c_{m,n,k}(i,j) u_i \otimes v_j$$

$$u_i \otimes v_j = \sum_k C_{m,n,k}(i,j) w_p^k$$

# Wigner's Formula (1931)

$$C'_{m,n,k}(i,j) = \sqrt{\frac{(m+n-2k+1)(m+n-i-j-k)!(i+j-k)!i!j!k!}{(m-i)!(n-j)!(m-k)!(n-k)!(m+n-k+1)!}} \cdot S$$

where

$$S = \sum_{s=0}^k (-1)^{i-s} \frac{(m-i+s)!(n-k+i-s)!}{s!(i+j-k-s)!(i-s)!(k-i+s)!}$$

Traditionally studied through representation theory, Jacobi polynomials and hypergeometric series of type  ${}_3F_2$

# Wigner's Formula (1931)

With binomial coefficients and shift  $s \rightarrow i - s$ ,

$$C'_{m,n,k}(i,j) = \sqrt{\frac{(m+n-2k+1) \binom{m+n-k}{m-i, n-j, *}}{(m+n-k+1) \binom{m+n-k}{m-k, n-k, k} \binom{m+n-k}{i, j, *}}} \cdot S$$

where

$$S = \sum_{s=0}^k (-1)^s \binom{i+j-k}{i-s} \binom{m-s}{k-s} \binom{n-k+s}{s}$$

## Racah's Formula (1942)

$$C'_{m,n,k}(i,j) = \sqrt{\frac{(m+n-2k+1) \binom{m+n-k}{m-k, n-k, k}}{(m+n-k+1) \binom{m+n-k}{m-i, n-j, *} \binom{m+n-k}{i, j, *}}} \cdot S'$$

where

$$S' = \sum_{s=0}^k (-1)^s \binom{k}{s} \binom{m-k}{i-s} \binom{n-k}{j-k+s}.$$

# Techniques:

## Traditional:

1. ONB and normalizations
2. Special functions and hypergeometric series
3. weighted recurrences

## Characteristic Zero:

1. no normalizations
2. Generating function and Binomial Identity
3. Pascal's recurrence

# Combinatorial Toolbox:

1. highest weight vectors with integral coordinates (LMC)
2. Clebsch-Gordan coefficients with rational values
3. elementary generating function
4. recurrence relations ( $144 = 2 \cdot 72$ )
5. Regge symmetries (72 symmetries, subgroup of  $S_6$ )
6. Elementary algorithm for bulk computation

# Open Problem:

When does  $C_{m,n,k}(i,j) = 0$ ?